

Outline of Lecture

- **Multiplication Algorithms**
- **Multiplication Hardware**

Multiplication

- Now we want to include multiplication in the design of our ALU - multiplication is much more complicated than addition or subtraction.
- More complicated than addition
 - accomplished via shifting and addition
- More time and more area

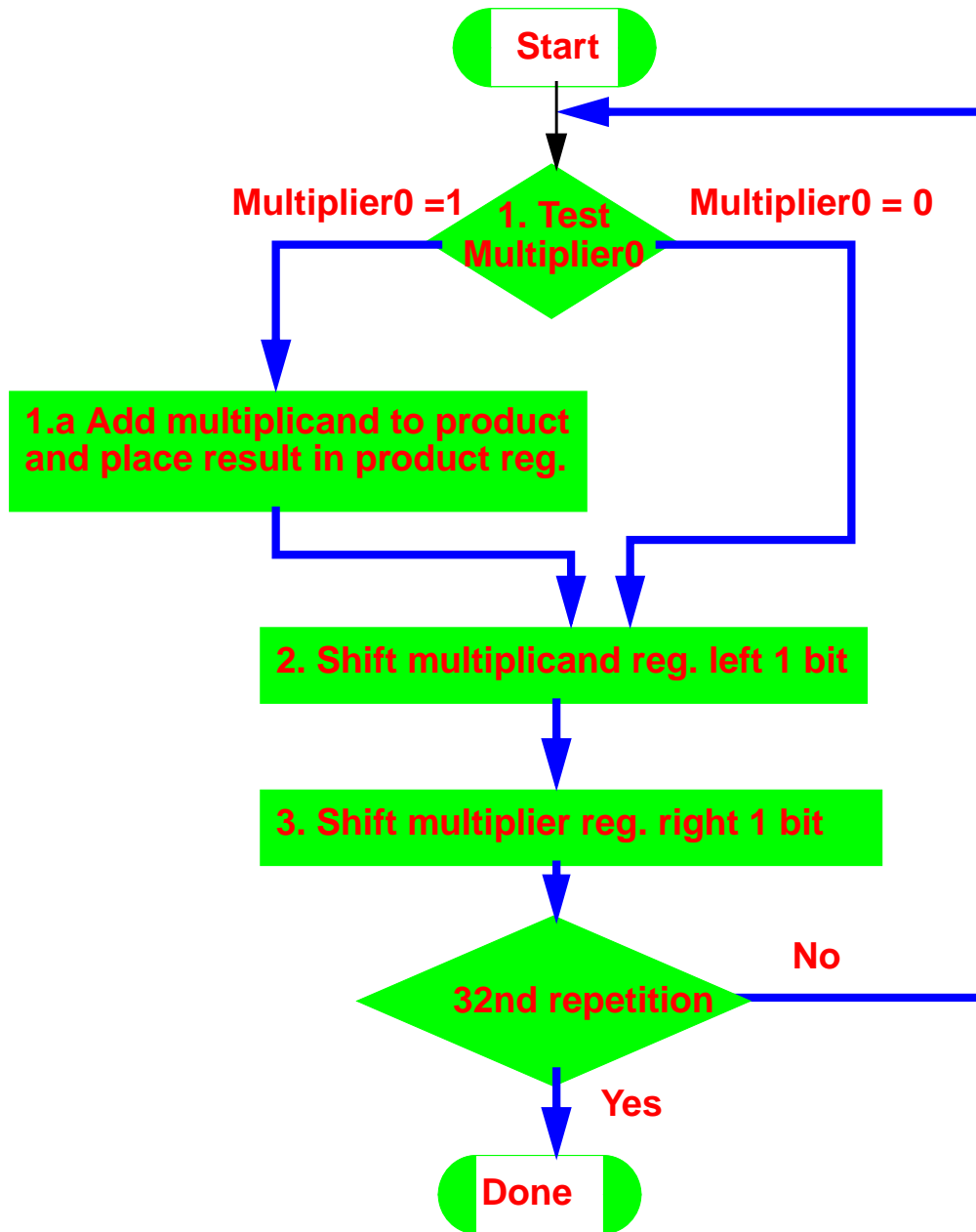
Simple Example

- **Paper-and-pencil example:**

Multiplicand	1 0 0 0
Multiplier	1 0 0 1
	<hr style="width: 100%;"/> 1 0 0 0 0 0 0 0 0 0 0 0 1 0 0 0 <hr style="width: 100%;"/> 1 0 0 1 0 0 0
Product	

- **m bits \times n bits = $m + n$ bit product**
- **if the digit is 0 \Rightarrow place (0 \times multiplicand)**
- **if the digit is 1 \Rightarrow place (1 \times multiplicand)**

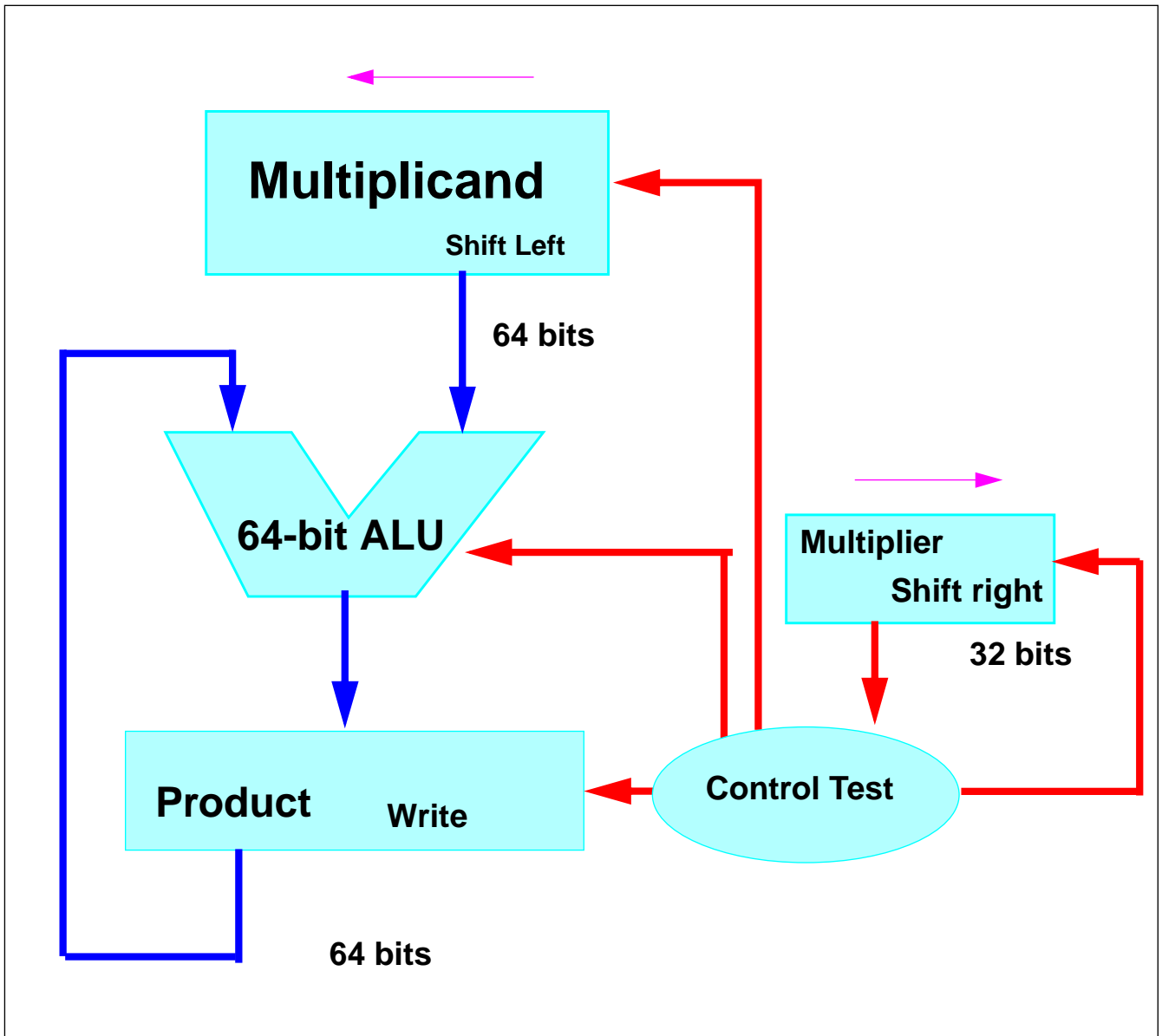
Multiplication



Multiplication Hardware

- The hardware for multiplication simply follows the flow of the above paper-and-pencil example.
- We use a 64-bit Multiplicand register, 64-bit ALU, 64-bit Product register, and a 32-bit Multiplier register.
 - The 32-bit multiplicand starts in the right half of the multiplicand register, and is shifted left 1 bit on each step.
 - The multiplier is shifted in the opposite direction at each step.
 - The product is initially 0.
 - The control decides when to shift the Multiplicand and Multiplier registers and when to write new values into the product register.

Multiplication Hardware



Example

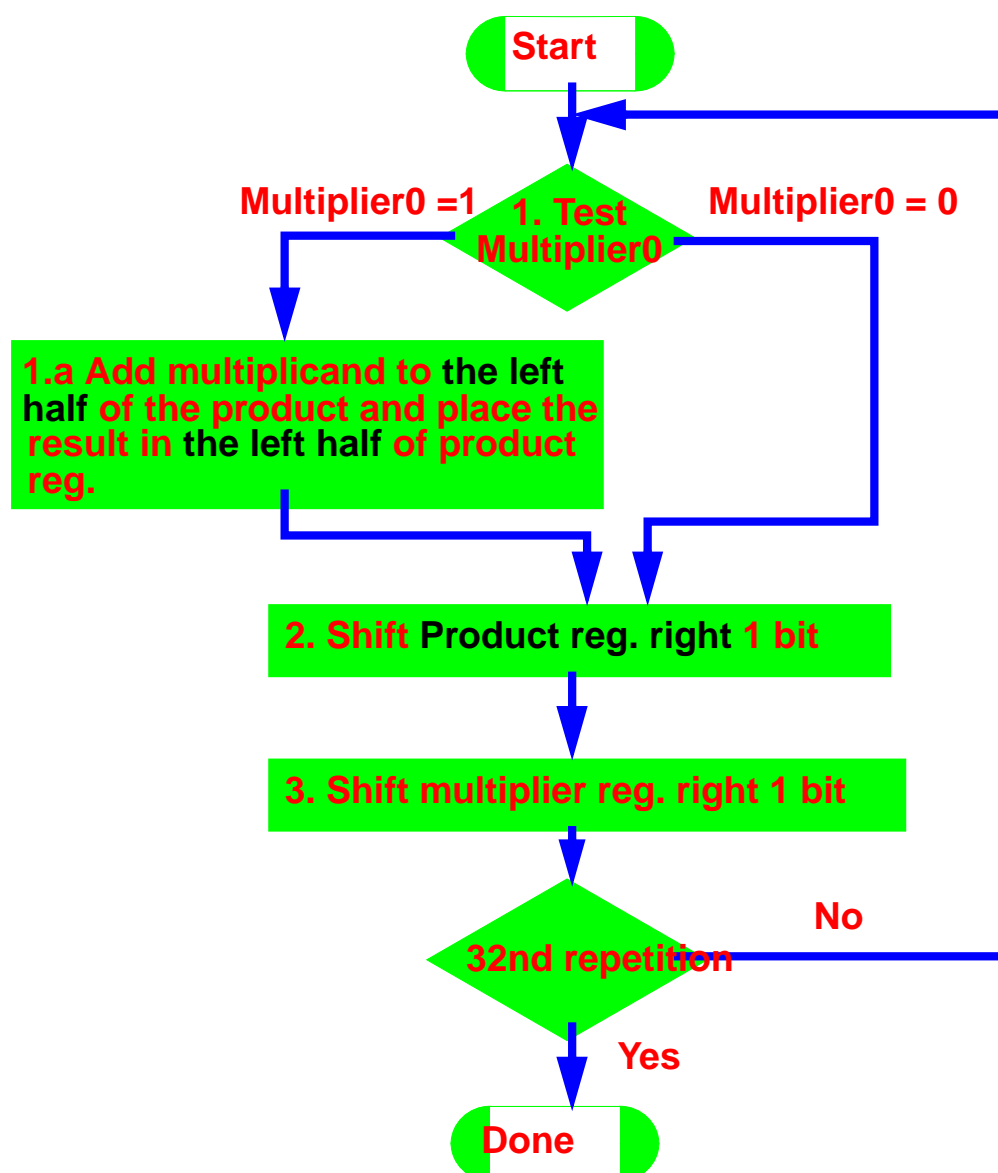
Using 4-bit numbers, multiply $0010_2 \times 0011_2$.

Solution

Iteration	Step	Multiplier	Multiplicand	Product
0	Initial Values	0011	0000 0010	0000 0000
1	1a: 1 \Rightarrow Prod=Prod+Mcand	0011	0000 0010	0000 0010
	2: Shift left multiplicand	0011	0000 0100	0000 0010
	3: Shift right multiplier	0001	0000 0100	0000 0010
2	1a: 1 \Rightarrow Prod=Prod+Mcand	0001	0000 0100	0000 0110
	2: Shift left multiplicand	0001	0000 1000	0000 0110
	3: Shift right multiplier	0000	0000 1000	0000 0110
3	1: 0 \Rightarrow no operation	0000	0000 1000	0000 0110
	2: Shift left multiplicand	0000	0001 0000	0000 0110
	3: Shift right multiplier	0000	0001 0000	0000 0110
4	1: 0 \Rightarrow no operation	0000	0001 0000	0000 0110
	2: Shift left multiplicand	0000	0001 0000	0000 0110
	3: Shift right multiplier	0000	0010 0000	0000 0110

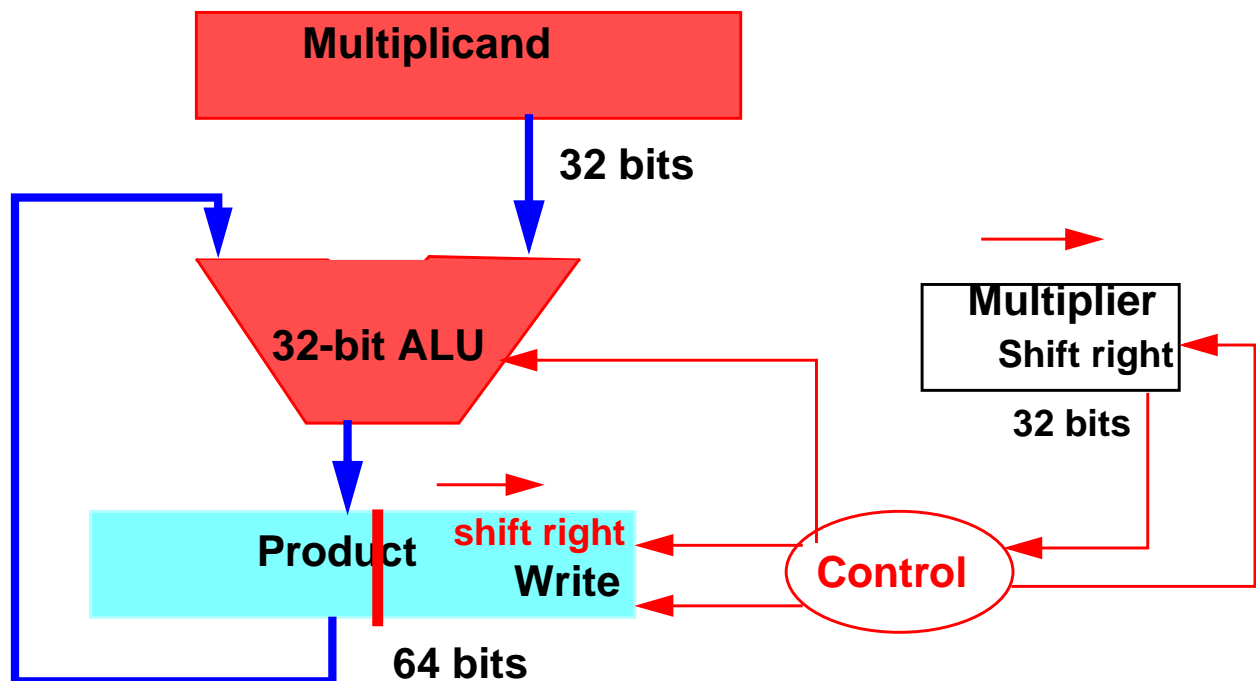
Second Version

- The hardware for the multiplication needs a 64-bit ALU. However, since half of the bits of the multiplicand are always 0 - we can achieve the same functionality by simply using a 32-bit ALU.



Second Version

- The hardware for this multiplication needs just 32-bit ALU - it can be implemented using our original MIPS ALU.



Signed Multiplication

- For signed multiplication: we simply convert the multiplier and multiplicand to positive numbers, and we remember the original signs. The algorithm will run for 31 iterations. After that, we restore the sign.

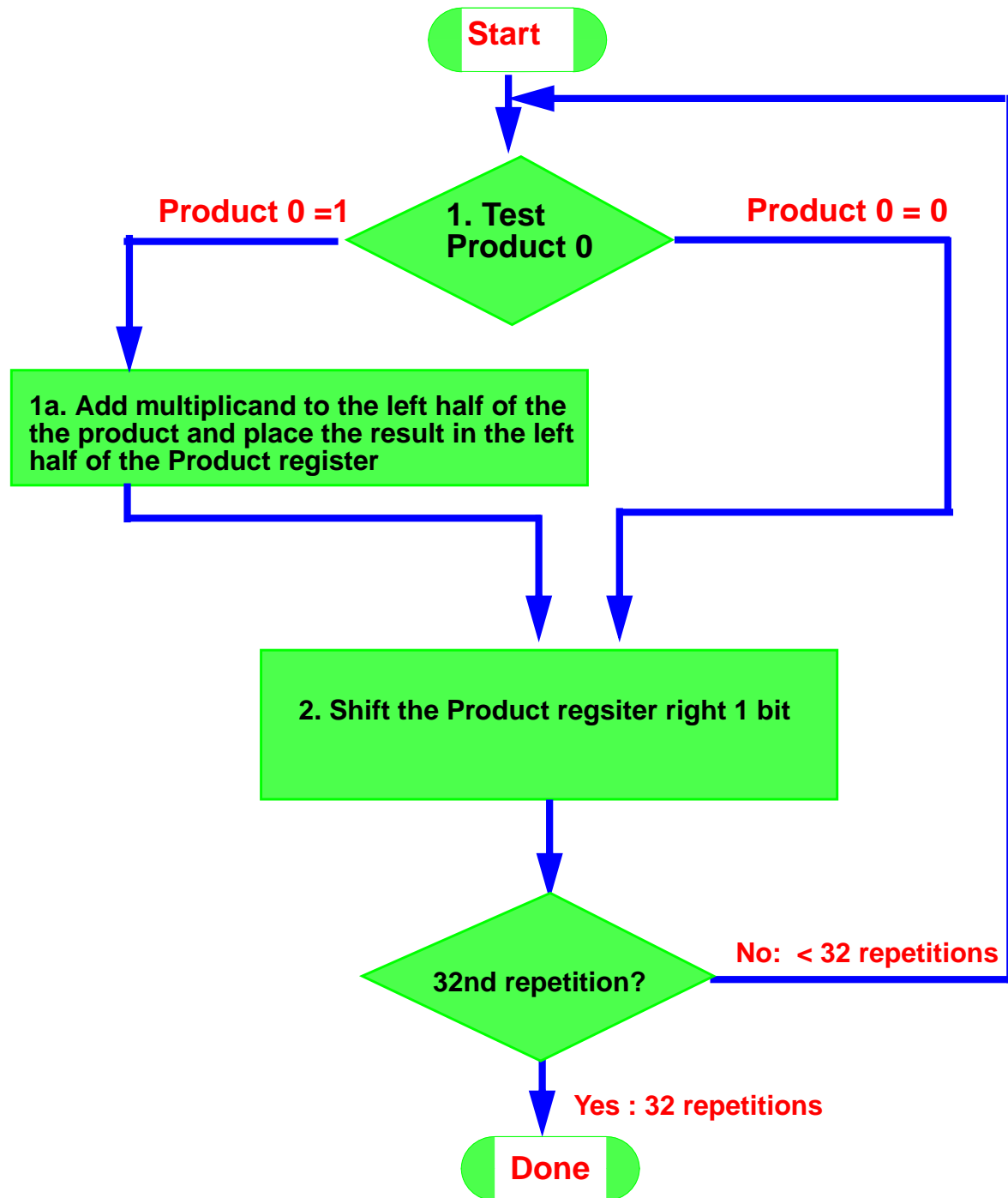
Example

Using 4-bit numbers, multiply $0010_2 \times 0011_2$.

Solution

Iteration	Step	Multiplier	Multiplicand	Product
0	Initial Values	0011	0010	0000 0000
1	1a: 1 \Rightarrow Prod=Prod+Mcand	0011	0010	0010 0000
	2: Shift right product	0011	0010	0001 0000
	3: Shift right multiplier	0001	0010	0001 0000
2	1a: 1 \Rightarrow Prod=Prod+Mcand	0001	0010	0011 0000
	2: Shift right product	0001	0010	0001 1000
	3: Shift right multiplier	0000	0010	0001 1000
3	1: 0 \Rightarrow no operation	0000	0010	0001 1000
	2: Shift right product	0000	0010	0000 1100
	3: Shift right multiplier	0000	0010	0000 1100
4	1: 0 \Rightarrow no operation	0000	0010	0000 1100
	2: Shift right product	0000	0010	0000 0110
	3: Shift right multiplier	0000	0010	0000 0110

Final Version



Final Version Hardware

